

Nonadditivity of Rains' bound for distillable entanglement

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Rains' bound is arguably the best known upper bound of the distillable entanglement by operations completely preserving positivity of partial transpose (PPT) and was conjectured to be additive and coincide with the asymptotic relative entropy of entanglement. We disprove both conjectures by explicitly constructing a special class of mixed two-qubit states. We then introduce an additive semidefinite programming lower bound (E_M) for the asymptotic Rains' bound, and it immediately becomes a computable lower bound for entanglement cost of bipartite states. Furthermore, E_M is also proved to be the best known upper bound of the PPT-assisted deterministic distillable entanglement and gives the asymptotic rates for all pure states and some class of genuinely mixed states.

I. INTRODUCTION

Entanglement plays a crucial role in quantum physics and is the key resource in many quantum information processing tasks. So it is quite natural and important to develop a theoretical framework to describe and quantify it. In spite of a series of remarkable recent progress in the theory of entanglement (for reviews see, e.g., Refs. [1–4]), many fundamental questions or challenges still remain open. One of the most significant questions is to determine the *distillable entanglement* E_D , i.e. the highest rate at which one can obtain maximally entangled states from an entangled state by local operations and classical communication (LOCC) [5, 6]. This fundamental measure fully captures the ability of given state shared between distant parties (Alice and Bob) to generate strongly correlated qubits in order to allow reliable quantum teleportation or quantum cryptography. However, up to now, how to calculate E_D for general quantum states still remains unknown. Also, in many practical applications, the resources are finite and the number of prepared states is limited. It is also of importance to study the deterministic distillable entanglement of finite entanglement transformations.

To evaluate the distillation rates efficiently, one possible way is to find computable upper bounds. A well-known upper bound of the distillable entanglement is the relative entropy of entanglement (REE) [7–9], which expresses the minimal distinguishability between the given state and all possible separable states. An improved bound is the Rains' bound [10], which is arguably the best known upper bound of distillable entanglement. The best known SDP upper bound is introduced in Ref. [11] and it is an improved version of the *logarithmic negativity* [12, 13]. In Ref [13], the logarithmic negativity is

proved to be a proper entanglement monotone for the first time. Other known upper bounds of E_D are studied in Refs. [8, 14–16]. Unfortunately, most of these known upper bounds are difficult to compute [17] and usually easily computable only for states with high symmetries, such as Werner states, isotropic states, or the family of “iso-Werner” states [5, 14, 18, 19].

Entanglement cost E_C [5, 20] is another fundamental measure in entanglement theory, which quantifies the rate for converting maximally entangled states to the given state by LOCC alone. However, computing E_C is NP-hard [17] and the entanglement cost is known only for a few of quantum states [21–23]. Even under the PPT operations, there are only bounds for the exact entanglement cost [24].

Since both distillable entanglement and the entanglement cost are important but difficult to compute, it is of great significance to find the best approach to efficiently evaluate them. As Rains' bound is proved to be equal to the asymptotic relative entropy of entanglement (with respect to PPT states) for Werner states [25] and orthogonally invariant states [26], one open problem is to determine whether these two quantities always coincide [1]. Another significant open problem is whether Rains' bound is additive, and it was conjectured in Ref. [26] that Rains' bound might be additive for arbitrary quantum states.

In this paper, we resolve the above two open problems about Rains' bound by explicitly exhibiting a special class of two-qubit states whose Rains' bound and relative entropy of entanglement are known. We show that the Rains' bound is not additive and thereby the asymptotic (or regularized) Rains' bound will give a better upper bound on distillable entanglement. Meanwhile, the asymptotic relative entropy of entanglement (w.r.t. PPT states) of these two-qubit states is strictly smaller than the Rains' bound. Furthermore, an SDP lower bound E_M for the asymptotic Rains' bound is introduced and it is the first computable lower bound for entanglement cost of general bipartite quantum states. Meanwhile, this

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bound is proved to be the best known upper bound of the deterministic distillable entanglement, which gives the PPT-assisted asymptotic rate for some states, including all the pure states and the mixed states $\rho^{(\alpha)}$ ($0 < \alpha \leq 0.2$) in Ref. [11].

Before we present our main results, let us first review some notations and preliminaries. In the following we will frequently use symbols such as A (or A') and B (or B') to denote (finite-dimensional) Hilbert spaces associated with Alice and Bob, respectively. The set of linear operators over A is denoted by $\mathcal{L}(A)$. Note that for a linear operator R over a Hilbert space, we define $|R| = \sqrt{R^\dagger R}$, and the trace norm of R is given by $\|R\|_1 = \text{Tr } |R|$, where R^\dagger is the conjugate transpose of R . The operator norm $\|R\|_\infty$ is defined as the maximum eigenvalue of $|R|$. A positive semidefinite operator $E_{AB} \in \mathcal{L}(A \otimes B)$ is said to be PPT if $E_{AB}^{T_B} \geq 0$, i.e., $(|i_A j_B\rangle\langle k_A l_B|)^{T_B} = |i_A l_B\rangle\langle k_A j_B|$.

The concise definition of entanglement of distillation by LOCC is given in Ref. [1] as follows:

$$E_D(\rho_{AB}) = \sup\{r : \lim_{n \rightarrow \infty} [\inf_{\Lambda} \|\Lambda(\rho_{AB}^{\otimes n}) - \Phi(2^{rn})\|_1] = 0\},$$

where Λ ranges over LOCC operations and $\Phi(d) = 1/d \sum_{i,j=1}^d |ii\rangle\langle jj|$ represents the standard $d \otimes d$ maximally entangled state. When Λ ranges over PPT operations, the PPT-assisted distillable entanglement is defined by $E_{D,PPT}$.

The Rains' bound was introduced in Ref. [10] and refined in Ref. [26] as a convex optimization problem as follows:

$$R(\rho) = \min S(\rho||\tau) \text{ s.t. } \tau \geq 0, \text{Tr } |\tau^{T_B}| \leq 1. \quad (1)$$

In this formula, $S(\rho||\sigma) = \text{Tr}(\rho \log \rho - \rho \log \sigma)$ denotes the relative Von Neumann entropy. Rains' bound is important in entanglement theory and the generalized Rains information of a quantum channel is recently proved to be a strong converse rate for quantum communication [32].

The relative entropy of entanglement (REE) [7–9] with respect to the PPT states is given by the following convex optimization problem:

$$E_{R,PPT}(\rho) = \min S(\rho||\sigma) \text{ s.t. } \sigma, \sigma^{T_B} \geq 0, \text{Tr}(\sigma) = 1. \quad (2)$$

For a general bipartite state ρ , it holds that $E_{R,PPT}(\rho) \geq R(\rho)$. However, $E_{R,PPT}(\rho)$ equals to $R(\rho)$ for every two-qubit state ρ [27] or the bipartite state with one qubit subsystem [28]. In particular, a two-qubit full-rank state σ is the closest separable state of any state ρ in the following form [27, 29]:

$$\rho = \sigma - xG(\sigma), \quad (3)$$

and

$$G(\sigma) = \sum_{i,j} G_{i,j} |v_i\rangle\langle v_i| (|\phi\rangle\langle\phi|)^{T_B} |v_j\rangle\langle v_j|, \quad (4)$$

with $\text{span}(|\phi\rangle)$ is the kernel (or null space) of σ^{T_B} and $G_{i,j} = \lambda_i$ when $\lambda_i = \lambda_j$ and $G_{i,j} = (\lambda_i - \lambda_j)/(\ln \lambda_i - \ln \lambda_j)$ when $\lambda_i \neq \lambda_j$, where λ_i and $|v_i\rangle$ are the eigenvalues and eigenvectors of σ , respectively.

The asymptotic relative entropy of entanglement is given by

$$E_{R,PPT}^\infty(\rho) = \inf_{n \geq 1} \frac{1}{n} E_{R,PPT}(\rho^{\otimes n}). \quad (5)$$

The numerical estimation of relative entropy of entanglement with respect to the PPT states is introduced in Refs. [30, 31], i.e., can be estimated by a Matlab program. Suppose that the estimation of $E_{R,PPT}(\rho)$ by in Refs. [30, 31] is $E_R^+(\rho)$, and the inequality $E_R^+(\rho) = S(\rho||\sigma) \geq E_{R,PPT}(\rho)$ holds since the algorithm indeed provides a feasible PPT state σ which is almost optimal. This algorithm is implemented by CVX [33] (a Matlab software for disciplined convex programming) and QETLAB [35]. In low dimensions, this algorithm provides an estimation $E_R^+(\rho)$ with an absolute error smaller than 10^{-3} , i.e. $E_{R,PPT}(\rho) + 10^{-3} \geq E_R^+(\rho) \geq E_{R,PPT}(\rho)$.

The SDP upper bound on distillable entanglement $E_W(\rho) = \log W(\rho)$ for a bipartite state ρ is introduced in Ref. [11], i.e.,

$$W(\rho) = \max \text{Tr } \rho R_{AB}, \quad |R_{AB}^{T_B}| \leq \mathbb{1}, R_{AB} \geq 0. \quad (6)$$

Semidefinite programming (SDP) [36] is a powerful tool in quantum information theory with many applications (e.g., [37–46]), which can be implemented by CVX [33] and QETLAB [35].

II. MAIN RESULTS

A. Nonadditivity of Rains' Bound

We first introduce a class of two-qubit states ρ_r whose closest separable states can be derived by the result in Ref. [27]. Thus, the Rains' bound of ρ_r is exactly given. Then we apply the algorithm in Refs. [30, 31] to demonstrate the gap between $R(\rho^{\otimes 2})$ and $R(\rho)$.

Theorem 1 *There exists a two-qubit state ρ such that*

$$R(\rho^{\otimes 2}) < 2R(\rho).$$

Meanwhile,

$$E_R^\infty(\rho) < R(\rho).$$

Proof Firstly, we construct two-qubit states ρ_r and σ_r satisfying Eq. (3). Then we have $R(\rho_r) = S(\rho_r||\sigma_r)$. Suppose that

$$\begin{aligned} \sigma_r = & \frac{1}{4} |00\rangle\langle 00| + \frac{1}{8} |11\rangle\langle 11| + r |01\rangle\langle 01| \\ & + \left(\frac{5}{8} - r\right) |10\rangle\langle 10| + \frac{1}{4\sqrt{2}} (|01\rangle\langle 10| + |10\rangle\langle 01|). \end{aligned}$$

The positivity of σ_r requires that $\frac{5-\sqrt{17}}{16} \leq r \leq \frac{5+\sqrt{17}}{16}$. Assume that $r \geq 5/8 - r$ and we can further choose $0.3125 \leq r \leq 0.57$ for simplicity.

Meanwhile, let us choose

$$\rho_r = \frac{1}{8}|00\rangle\langle 00| + x|01\rangle\langle 01| + \frac{7-8x}{8}|10\rangle\langle 10| + \frac{32r^2 - (6+32x)r + 10x + 1}{4\sqrt{2}}(|01\rangle\langle 10| + |10\rangle\langle 01|)$$

with

$$x = r + \frac{32r^2 - 10r + 1}{256r^2 - 160r + 33} + \frac{(16r - 5)y^{-1}}{32 \ln(5/8 - y) - 32 \ln(5/8 + y)}$$

and $y = (4r^2 - 5r/2 + 33/64)^{1/2}$. It is clear that $\text{Tr } \rho_r = 1$ and we set $0.3125 \leq r \leq 0.5480$ to ensure the positivity of ρ_r .

One can readily verify that $\rho_r = \sigma_r - 3G(\sigma_r)/2$. Therefore, σ_r is the closest separable state (CSS) for ρ_r and we have that

$$R(\rho_r) = E_{R,PPT}(\rho_r) = S(\rho_r || \sigma_r). \quad (7)$$

In particular, let us first choose $r_0 = 0.547$, the Rains' bound of ρ_{r_0} is given by

$$R(\rho_{r_0}) = E_{R,PPT}(\rho_{r_0}) = S(\rho_{r_0} || \sigma_{r_0}) \simeq 0.3891999.$$

Furthermore, applying the algorithm in Refs. [30, 31], we can find a PPT state σ_0 such that

$$E_R^+(\rho_{r_0}^{\otimes 2}) = S(\rho_{r_0}^{\otimes 2} || \sigma_0) \simeq 0.7683307.$$

The numerical value of relative entropy here is calculated based on the Matlab function “logm” [34] and the function “Entropy” in QETLAB [35]. In this case, the accuracy is guaranteed by the fact $\|e^{\log m(\sigma_{r_0})} - \sigma_{r_0}\|_1 \leq 10^{-16}$ and $\|e^{\log m(\sigma_0)} - \sigma_0\|_1 \leq 10^{-14}$. Noting that the difference between $2R(\rho_{r_0})$ and $E_R^+(\rho_{r_0}^{\otimes 2})$ is already 1.00691×10^{-2} , we have that

$$R(\rho_{r_0}^{\otimes 2}) \leq E_{R,PPT}(\rho_{r_0}^{\otimes 2}) \leq E_R^+(\rho_{r_0}^{\otimes 2}) < 2R(\rho_{r_0}).$$

It is also easy to observe that

$$E_{R,PPT}^\infty(\rho_{r_0}) \leq \frac{1}{2}E_{R,PPT}(\rho_{r_0}^{\otimes 2}) < R(\rho_{r_0}).$$

When $0.45 \leq r \leq 0.548$, we show the gap between $2R(\rho_r)$ and $E_R^+(\rho_r^{\otimes 2})$ in FIG. 1. \square

Since Rains' bound is not additive, the asymptotic Rains' bound [47] can provide better upper bound on the distillable entanglement, i.e.,

$$E_{D,PPT}(\rho) \leq R^\infty(\rho) = \inf_{n \geq 1} \frac{1}{n} R(\rho^{\otimes n}) \leq R(\rho), \quad (8)$$

and the last inequality can be strict.

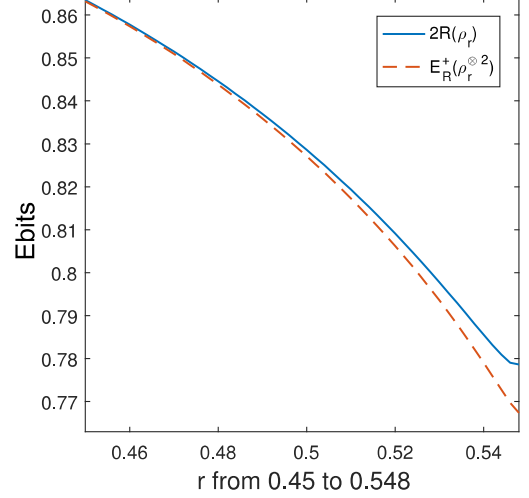


FIG. 1: This plot demonstrates the difference between $2R(\rho_r)$ and $E_R^+(\rho_r^{\otimes 2})$ for $0.45 \leq r \leq 0.548$. The dashed line depicts $E_R^+(\rho_r^{\otimes 2})$ while the solid line depicts $2R(\rho_r)$.

B. A SDP lower bound for entanglement cost

Since computing the entanglement cost of a bipartite state is very difficult, we introduce an efficiently computable lower bound to evaluate the entanglement cost.

For a bipartite quantum state ρ , we introduce

$$E_M(\rho) = -\log M(\rho) = -\log \max \text{Tr } P_{AB} V_{AB}, \quad (9)$$

s.t. $\text{Tr } |V_{AB}^{T_B}| = 1, V_{AB} \geq 0,$

where P_{AB} is the projection onto the support of ρ . And $M(\rho)$ is also given by the following SDP:

$$M(\rho) = \max \text{Tr } P_{AB} Z_{AB}, \quad (10)$$

s.t. $\text{Tr}(X_{AB} + Y_{AB}) = 1,$
 $Z_{AB} \leq (X_{AB} - Y_{AB})^{T_B}$
 $X_{AB}, Y_{AB}, Z_{AB} \geq 0,$

And its dual SDP is given by

$$M(\rho) = \min \|R_{AB}^{T_B}\|_\infty, \text{ s.t. } R_{AB} \geq P_{AB}. \quad (11)$$

The optimal values of the primal and the dual SDPs above coincide by strong duality, which can be proved by Slaters theorem.

For any two bipartite states ρ_{AB} and $\sigma_{A'B'}$, by utilizing semidefinite programming duality, it is not difficult to prove that

$$E_M(\rho_{AB} \otimes \sigma_{A'B'}) = E_M(\rho_{AB}) + E_M(\sigma_{A'B'}).$$

Furthermore, for any state bipartite ρ , $E_M(\rho) = 0$ if and only if $\text{supp}(\rho)$ contains the support of a PPT state σ , i.e. $\text{supp}(\sigma) \subseteq \text{supp}(\rho)$. Too see this, if there exists PPT

state σ such that $\text{supp}(\sigma) \subseteq \text{supp}(\rho)$, then $E_M(\rho) = 0$. On the other hand, if any state σ satisfies $\text{supp}(\sigma) \subseteq \text{supp}(\rho)$ is NPPT. Let the optimal solution to SDP (9) be V , where $V \geq 0$ and $\text{Tr}|V^{T_B}| = 1$. It is clear that $\text{Tr} V \leq 1$. Thus, we have $\text{Tr} V = 1$ when $E_M(\rho) = 0$. Hence, V is a PPT state and $\text{supp}(V) \subseteq \text{supp}(\rho)$. This leads to a contradiction.

Theorem 2 For any bipartite state ρ ,

$$E_M(\rho) \leq R^\infty(\rho) \leq E_C(\rho).$$

Proof Suppose that the optimal solution to Eq. (1) of $R(\rho)$ is V , then $V \geq 0$ and $\text{Tr}|V^{T_B}| = t \leq 1$. Thus, V/t is a feasible solution to SDP (9) of $M(\rho)$, this means that $E_M \leq -\log \text{Tr} PV/t = -\log \text{Tr} PV + \log t$, where P is the projection onto $\text{supp}(\rho)$.

On the other hand, let $\mathcal{N}(\sigma) = P\sigma P + (\mathbb{1} - P)\sigma(\mathbb{1} - P)$, then by the monotonicity of quantum relative entropy,

$$\begin{aligned} S(\rho||V) &\geq S(\mathcal{N}(\rho)||\mathcal{N}(V)) = S(\rho||PVP) \\ &= S(\rho||\frac{PVP}{\text{Tr} PVP}) - \log \text{Tr} PV \\ &\geq -\log \text{Tr} PV \geq E_M(\rho). \end{aligned} \quad (12)$$

Noting that $E_M(\cdot)$ is additive, we have that

$$E_M(\rho) \leq \inf_{n \geq 1} \frac{1}{n} R(\rho^{\otimes n}) = R^\infty(\rho).$$

Finally, it is clear that

$$E_M(\rho) \leq R^\infty(\rho) \leq E_{R,PPT}^\infty(\rho) \leq E_C(\rho),$$

where the last inequality is from Ref. [47]. \square

Remark As an application of this lower bound, one can also give an SDP lower bound for the *entanglement cost of quantum channels* [48], i.e. the rate of entanglement (ebits) needed to asymptotically simulate a quantum channel \mathcal{N} with free classical communication. \square

C. Deterministic distillable entanglement

In this section, we show that E_M is the best upper bound on the deterministic distillable entanglement of bipartite states. The bipartite pure state case is completely solved in Refs. [49, 50]. For a general state, the PPT-assisted deterministic distillation rates depend only on the support of this state [11]. Note that the support $\text{supp}(\rho)$ of a state ρ is defined to be the space spanned by the eigenvectors with non-zero eigenvalues of ρ . The exact value of the one-copy PPT-assisted deterministic distillation rate of a given bipartite state ρ is $E_{0,D,PPT}^{(1)}(\rho) = -\log W_0(\rho)$ [11], where $W_0(\rho)$ is given by

$$W_0(\rho) = \min \|R^{T_B}\|_\infty, \text{ s.t. } P_{AB} \leq R \leq \mathbb{1}_{AB}. \quad (13)$$

Here, P_{AB} is the projection onto $\text{supp}(\rho)$.

Theorem 3 For any bipartite state ρ ,

$$E_{0,D,PPT}(\rho) \leq E_M(\rho) \leq E_W(\rho).$$

Proof To prove $E_{0,D,PPT}(\rho) \leq -\log M(\rho)$, suppose that the optimal solution to SDP (13) of $W_0(\rho)$ is R_0 . It is clear that R_0 is also a feasible solution to SDP (11) of $M(\rho)$. Thus, $W_0(\rho) = \|R_0^{T_B}\|_\infty \geq M(\rho)$. Furthermore, $W_0(\rho^{\otimes n}) \geq M(\rho^{\otimes n}) = M(\rho)^n$.

Hence,

$$\begin{aligned} E_{0,D,PPT}(\rho) &= \lim_{n \rightarrow \infty} -\frac{1}{n} \log W_0(\rho^{\otimes n}) \\ &\leq \lim_{n \rightarrow \infty} -\frac{1}{n} \log M(\rho)^n = E_M(\rho). \end{aligned}$$

Finally, to prove $E_M(\rho) \leq E_W(\rho)$, suppose that the optimal solution to SDP (11) is R , then we have $R \geq P \geq 0$. Let $R' = R/\|R^{T_B}\|_\infty$ and it is easy to see the positivity of R' and the fact that $|R'^{T_B}| \leq \mathbb{1}$, which means that R' is a feasible solution to SDP (6). Therefore, $E_W(\rho) \geq \log \text{Tr} \rho R' \geq \log \text{Tr} \rho P / \|R^{T_B}\|_\infty = -\log \|R^{T_B}\|_\infty = E_M(\rho)$. \square

Remark For any bipartite state ρ , if the support of ρ contains a PPT state σ , then $E_M(\rho) = 0$ and we have that $E_{0,D,PPT}(\rho) = 0$. Thus ρ is bound entanglement for exact distillation under both LOCC or PPT operations. \square

We further show the estimation of Theorem 3 in Fig.2 by a class of $3 \otimes 3$ states in Ref. [11] defined by

$$\rho^{(\alpha)} = \frac{1}{3} \sum_{m=0}^2 (X^\dagger \otimes X)^m |\psi_0\rangle\langle\psi_0| (X \otimes X^\dagger)^m,$$

where $|\psi_0\rangle = \sqrt{\alpha}|00\rangle + \sqrt{1-\alpha}|11\rangle$ ($0 < \alpha \leq 0.5$) and $X = \sum_{j=0}^2 |j \oplus 1\rangle\langle j|$. An interesting fact is that $E_M(\rho^{(\alpha)})$ is tight for $E_{0,D,PPT}(\rho^{(\alpha)})$ when $0 < \alpha \leq 1/5$, which is proved in the following Proposition.

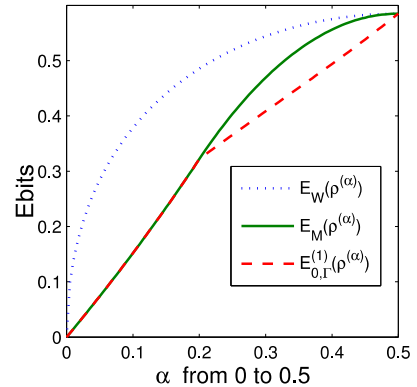


FIG. 2: This plot presents the estimation of $E_{D,PPT}(\rho^{(\alpha)})$ and $E_{0,D,PPT}(\rho^{(\alpha)})$. The dot line depicts $E_W(\rho^{(\alpha)})$, the dash line depicts $E_{0,D,PPT}^{(1)}(\rho^{(\alpha)})$ and the solid line depicts $E_M(\rho^{(\alpha)})$.

Proposition 4 For any bipartite state ρ with support projection P , suppose that the eigenvector $|\psi\rangle$ of P^{T_B} with the eigenvalue $\|P^{T_B}\|_\infty$ is a product state, then

$$E_{0,D,PPT}(\rho) = E_M(\rho) = -\log \|P^{T_B}\|_\infty \leq E_{D,PPT}(\rho).$$

Proof In Ref. [11], it shows that $E_{0,D,PPT}(\rho) \geq -\log \|P^{T_B}\|_\infty$. If $|\psi\rangle\langle\psi|$ is PPT, then we can choose $V = |\psi\rangle\langle\psi|$ and it is easy to see V is a feasible solution to SDP (9) of $M(\rho)$. Thus, $E_M(\rho) \leq -\log \text{Tr } P^{T_B} |\psi\rangle\langle\psi| = -\log \|P^{T_B}\|_\infty$. \square

For any pure state $|\phi\rangle\langle\phi|$, suppose that $|\phi\rangle$ has the Schmidt decomposition $|\phi\rangle = \sum_{i=1}^m \lambda_i |ii\rangle$ with $\lambda_1^2 \geq \dots \geq \lambda_m^2$ and $\sum_{i=1}^m \lambda_i^2 = 1$. Then $|\phi\rangle\langle\phi|^{T_B} = \sum_{i=1}^m \lambda_i^2 |ii\rangle\langle ii| + \sum_{i \neq j} \lambda_i \lambda_j |ji\rangle\langle ij|$. Thus, $\|P^{T_B}\|_\infty = \lambda_1^2$ and the corresponding eigenvector is $|11\rangle\langle 11|$. Hence, by Proposition 4, $E_{0,D,PPT}(|\phi\rangle\langle\phi|) = E_M(|\phi\rangle\langle\phi|) = -\log \| |\phi\rangle\langle\phi|^{T_B} \|_\infty$. This rate can be achieved by LOCC [50].

Example For the $\rho^{(\alpha)}$, when $0 < \alpha \leq 1/5$, we have that $E_{0,D,PPT}(\rho^{(\alpha)}) = E_M(\rho^{(\alpha)}) = -\log(1 - \alpha)$. Let $U = X^\dagger \otimes X$, the projection onto $\text{supp}(\rho^{(\alpha)})$ is $P_\alpha = \sum_{m=0}^2 U^m |\psi_0\rangle\langle\psi_0| (U^\dagger)^m$. Therefore,

$$P_\alpha^{T_B} = 2\sqrt{\alpha(1-\alpha)} |v_1\rangle\langle v_1| - \sqrt{\alpha(1-\alpha)} (|v_2\rangle\langle v_2| + |v_3\rangle\langle v_3|) + \sum_{m=0}^2 U^m [(1-\alpha)|11\rangle\langle 11| + \alpha|00\rangle\langle 00|] (U^\dagger)^m,$$

where $|v_1\rangle = \frac{1}{\sqrt{3}}(|01\rangle + |10\rangle + |22\rangle)$, $|v_2\rangle = \frac{1}{\sqrt{6}}|01\rangle + \frac{1}{\sqrt{6}}|10\rangle - \sqrt{\frac{2}{3}}|22\rangle$ and $|v_3\rangle = \frac{1}{\sqrt{2}}(|01\rangle - |10\rangle)$. When $0 < \alpha \leq 1/5$, we always have $1 - \alpha \geq 2\sqrt{\alpha(1-\alpha)}$. Therefore,

$\|P_\alpha^{T_B}\|_\infty = 1 - \alpha$ and the corresponding eigenvector is $|11\rangle\langle 11|$. Applying Proposition 4, the proof is done. \square

III. CONCLUSIONS AND DISCUSSIONS

We show that the Rains' bound is neither additive nor equal to the asymptotic relative entropy of entanglement by explicitly constructing a special class of mixed two-qubit states. We also show an SDP lower bound E_M for the asymptotic Rains' bound. These results solve two open problems in entanglement theory and provide an efficiently computable lower bound for the entanglement cost of general bipartite states. This bound also has desirable properties such as additivity under tensor product and vanishing if and only if the support of the given state contains some PPT state. We further show that E_M is the best upper bound for the deterministic distillable entanglement, which also gives the PPT-assisted deterministic distillation rate in some conditions, including all the pure states and some classes of the mixed states.

It is of great interest to determine whether the asymptotic Rains' bound and the PPT distillable entanglement always coincide. It would also be interesting to decide whether $E_{0,D,PPT}(\rho) = E_M(\rho)$ for any bipartite state ρ and to study the relationship between E_M and the newly established SDP lower bound of the PPT-assisted entanglement cost [51].

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